

Many-valued modal logic over a semi-primal algebra

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Joint work (in progress) with Alexander Kurz and Bruno Teheux

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Semi-primal algebras

The two-element Boolean algebra $\mathbf{2}$ is *primal*. That is, every map $f : 2^n \rightarrow 2$ is term-definable.

Definition (Foster, Pixley 1964)

A finite algebra L is *semi-primal* if every map $f : L^n \rightarrow L$ which preserves subalgebras is term-definable.

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In a *discriminator algebra* the ternary discriminator is term-definable:

$$t(x, y, z) = \begin{cases} z & \text{if } x = y \\ x & \text{if } x \neq y \end{cases}$$

Proposition (Pixley 1971)

A finite algebra \mathbf{L} is semi-primal if and only if it is a discriminator algebra and it has no internal isomorphisms.

In this talk we exclusively consider semi-primal algebras \mathbf{L} with an underlying bounded lattice structure $\mathbf{L}^b = (L, \wedge, \vee, 0, 1)$.

Proposition

Let \mathbf{L} be a finite algebra with bounded lattice reduct. Then \mathbf{L} is semi-primal if and only if for every $a \in L$ the map $T_a : L \rightarrow L$:

$$T_a(x) = \begin{cases} 1 & \text{if } x = a \\ 0 & \text{if } x \neq a \end{cases}$$

is term-definable.

Examples

- Chain-based: The $(n + 1)$ -element Łukasiewicz-chain

$$\mathbf{L}_n = (\{0, \frac{1}{n}, \dots, \frac{n-1}{n}, 1\}, \wedge, \vee, \oplus, \neg, 0, 1)$$

where

$$x \oplus y = \min(x + y, 1) \text{ and } \neg x = 1 - x.$$

Examples

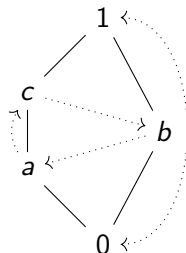
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- Not chain-based: $\mathbf{L} = (N_5, \wedge, \vee, ', 0, 1)$:



(Davey, Schumann, Werner 1991)

Stone duality and semi-primal natural duality

For $\mathbf{BA} = \mathbf{ISP}(\mathbf{2})$ we have the famous *Stone duality*:

$$\text{Stone} \begin{array}{c} \xrightarrow{\quad \Pi \quad} \\ \xleftarrow{\quad \Sigma \quad} \end{array} \mathbf{BA}$$

$\Sigma(\mathbf{B}) = \mathbf{BA}(\mathbf{B}, \mathbf{2})$ and $\Pi(X) = \text{Stone}(X, \mathbf{2})$.

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For $\mathcal{A} := \mathbf{ISP}(\mathbf{L})$ we have the similar *natural duality*:

$$\text{Stone}_{\mathbf{L}} \begin{array}{c} \xrightarrow{P} \\ \xleftarrow{S} \end{array} \mathcal{A}$$

“ $S(\mathbf{A}) = \mathcal{A}(\mathbf{A}, \mathbf{L})$ and $P(X) = \text{Stone}_{\mathbf{L}}(X, \mathbf{L})$ ”
(Keimel, Werner 1973)

The category Stone_L

The category Stone_L has objects (X, ν) where

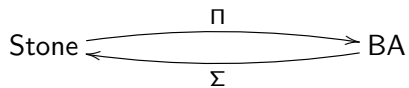
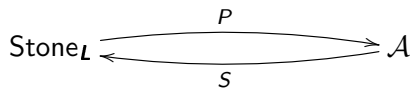
- $X \in \text{Stone}$,
- $\nu : X \rightarrow \mathbb{S}(L)$,
- $\nu^{-1}(\mathbf{S}\downarrow)$ is closed for every subalgebra $\mathbf{S} \leq L$.

A morphism $f : (X, \nu) \rightarrow (Y, \nu')$ is a continuous map with

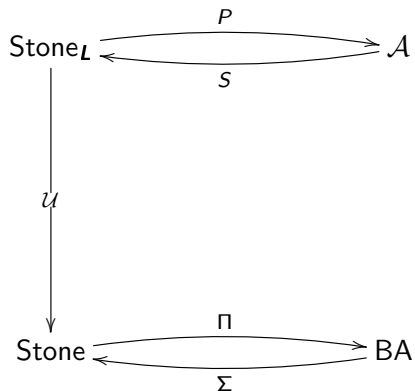
$$\nu'(f(x)) \subseteq \nu(x)$$

for every $x \in X$.

Relating the dualities

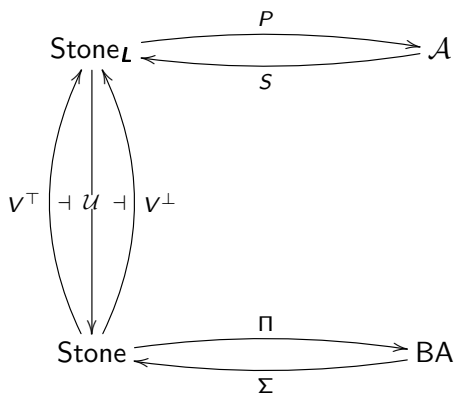


Relating the dualities



Let \mathcal{U} be the forgetful functor.

Relating the dualities

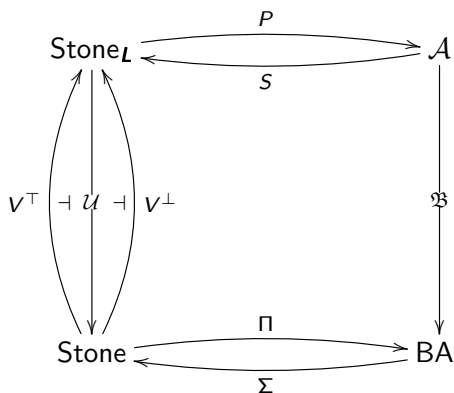


Let \mathcal{U} be the forgetful functor.

$V^\top(X) = (X, v^\top)$ where $v^\top(x) = L$ for all $x \in X$.

$V^\perp(X) = (X, v^\perp)$ where $v^\perp(x) = 2$ for all $x \in X$.

Relating the dualities



Let \mathcal{U} be the forgetful functor.

$V^\top(X) = (X, v^\top)$ where $v^\top(x) = L$ for all $x \in X$.

$V^\perp(X) = (X, v^\perp)$ where $v^\perp(x) = \langle 0, 1 \rangle$ for all $x \in X$.

The Boolean skeleton

Recall: For every $a \in L$ the map $T_a : L \rightarrow L$ given by

$$T_a(x) = \begin{cases} 1 & \text{if } x = a \\ 0 & \text{if } x \neq a \end{cases} \text{ is term-definable in } L.$$

Definition (Maruyama 2011)

Given any $\mathbf{A} \in \mathcal{A}$ we define $\mathfrak{B}(\mathbf{A}) = \{a \in A \mid T_1(a) = a\}$. The *Boolean skeleton* of \mathbf{A} is the Boolean algebra

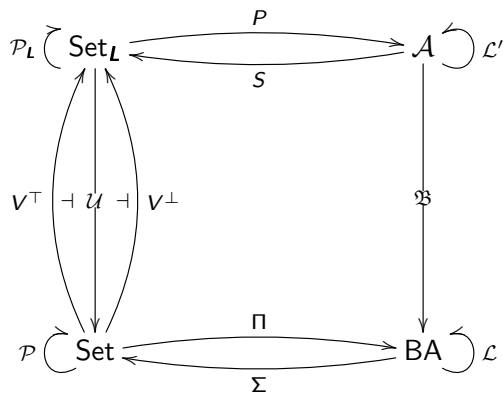
$$\mathfrak{B}(\mathbf{A}) = (\mathfrak{B}(\mathbf{A}), \wedge, \vee, T_0, 0, 1).$$

Theorem

Let $\mathbf{A} \in \mathcal{A}$. Then $h \mapsto h \upharpoonright_{\mathfrak{B}(\mathbf{A})}$ is a homeomorphism

$$\mathcal{A}(\mathbf{A}, L) \simeq \text{BA}(\mathfrak{B}(\mathbf{A}), 2).$$

Going modal, coalgebraically



\mathcal{P} -coalgebras: (Kripke-)frames

\mathcal{L} -algebras: Classical modal logic

\mathcal{P}' -coalgebras: L -(Kripke-)frames

\mathcal{L}' -algebras: Modal logic over L

Definition

An L -frame is a triple (W, R, v) such that

- (W, R) is a Kripke-frame,
- $(W, v) \in \text{Set}_L$, (that is, $v : W \rightarrow \mathbb{S}(L)$) and
- Compatibility: $wRw' \Rightarrow v(w') \subseteq v(w)$.

An L -model is given by $\mathfrak{M} = (W, R, v, Val)$ such that (W, R, v) is an L -frame and

$$Val : W \times \text{Prop} \rightarrow L$$

satisfies for all $w \in W, p \in \text{Prop}$:

$$Val(w, p) \in v(w).$$

We can extend Val to all modal formulas and define the validity relation $\mathfrak{M}, w \models \varphi$ iff $Val(w, \varphi) = 1$.

The ' L -versions' of the Goldblatt-Thomason theorem

The functor \mathfrak{B} - Boolean skeleton - is used in the proof of the following:

Theorem

Let F be a class of L -frames closed under ultrapowers. Then F is definable if and only if F is closed under disjoint unions, generated subalgebras and bounded morphic images, and it reflects canonical extensions.

The dually adjoint functor V^T is then used in the proof of the following:

Corollary

Let F be a class of Kripke-frames closed under ultrapowers. Then F is L -definable if and only if F is closed under disjoint unions, generated subalgebras and bounded morphic images, and it reflects ultrafilter extensions.

Thanks for your attention!