

Change of Base using Arrow Categories

Michael Winter

Department of Computer Science,
Brock University,
St. Catharines, Ontario, Canada, L2S 3A1
mwinter@brocku.ca

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Motivation

We use arrow categories to formulate fuzzy concepts and constructions within the language of abstract categories of relations.

- 1 Arrow categories form an abstract theory of L -valued relation where L is a complete Heyting algebra.
- 2 We present an algebraic framework for handling a change of base for L -relations, i.e., a functor from \mathcal{A}_1 to \mathcal{A}_2 .
- 3 Important to obtain a canonical embedding L -relations into type-2 L -relations, i.e., relations with membership values from $L \rightarrow L$, which is essential in the construction of type-2 fuzzy controllers.
 - The arrow category of type-2 L -relations is given by the Kleisli category of a (relational) monad based on the product functor.
 - The embedding maps an element u of L to the constant function $\bar{u} : L \rightarrow L$ defined by $\bar{u}(x) = u$.



Dedekind Categories

Definition

A Dedekind category \mathcal{R} is a category satisfying the following:

- 1 For all objects A and B the collection $\mathcal{R}[A, B]$ is a complete distributive lattice with operations $\sqcap, \sqcup, \sqsubseteq, \perp_{AB}, \top_{AB}$.
- 2 There is a monotone operation $\bar{}$ (called conversion) so that for all relations $Q : A \rightarrow B$ and $R : B \rightarrow C$:

$$(Q; R)^{\bar{}} = R^{\bar{}}; Q^{\bar{}}, \quad (Q^{\bar{}})^{\bar{}} = Q.$$

- 3 For all relations $Q : A \rightarrow B, R : B \rightarrow C$ and $S : A \rightarrow C$ the modular law holds:

$$Q; R \sqcap S \sqsubseteq Q; (R \sqcap Q^{\bar{}}; S).$$

- 4 For all relations $R : B \rightarrow C$ and $S : A \rightarrow C$ there is a relation $S/R : A \rightarrow B$ (called the left residual of S and R) so that for all $Q : A \rightarrow B$ the following holds:

$$Q; R \sqsubseteq S \iff Q \sqsubseteq S/R.$$



Special L -fuzzy Relations I

Relations corresponding to elements in L :



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scalar relations

Definition

- 1 A relation $J : A \rightarrow B$ is called an ideal iff $\pi_{AA}; J; \pi_{BB} = J$.
- 2 A relation $\alpha : A \rightarrow A$ is called a scalar on A iff $\alpha \sqsubseteq \mathbb{I}_A$ and $\pi_{AA}; \alpha = \alpha; \pi_{AA}$.



Arrow Categories I

Definition

An arrow category \mathcal{A} is a non-trivial Dedekind category ($\top_{AB} \neq \perp_{AB}$ for all objects A and B) together with two operations \uparrow and \downarrow satisfying the following:

- ① $R^\uparrow, R^\downarrow : A \rightarrow B$ for all $R : A \rightarrow B$.
- ② (\uparrow, \downarrow) is a Galois correspondence.
- ③ $(R^\sim; S^\downarrow)^\uparrow = R^{\uparrow\sim}; S^\downarrow$ for all $R : B \rightarrow A$ and $S : B \rightarrow C$.
- ④ $(Q \sqcap R^\downarrow)^\uparrow = Q^\uparrow \sqcap R^\downarrow$ for all $Q, R : A \rightarrow B$.
- ⑤ If $\alpha_A \neq \perp_{AA}$ is a non-zero scalar, then $\alpha_A^\uparrow = \mathbb{I}_A$.



Arrow Categories II

Theorem

Let \mathcal{A} be an arrow category. Then we have:

- 1 \mathcal{A} is uniform, i.e., we have $\pi_{AB}; \pi_{BC} = \pi_{AC}$ for all objects A, B and C .



Arrow Categories II

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Let \mathcal{A} be an arrow category. Then we have:

- 1 \mathcal{A} is uniform, i.e., we have $\pi_{AB}; \pi_{BC} = \pi_{AC}$ for all objects A, B and C .
- 2 The complete Heyting algebras $\text{Sc}(A)$ and $\text{Sc}(B)$ of scalar relations on any two objects A and B are isomorphic. We denote this lattice by $\text{Sc}[\mathcal{A}]$.



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- ② The complete Heyting algebras $\text{Sc}(A)$ and $\text{Sc}(B)$ of scalar relations on any two objects A and B are isomorphic. We denote this lattice by $\text{Sc}[\mathcal{A}]$.
- ③ The substructure \mathcal{A}^\downarrow of crisp relations, i.e., all relations R with $R^\downarrow = R$ (or equivalently $R^\uparrow = R$), is a Dedekind category satisfying the so-called Tarski rule, i.e., the equivalence $R \neq \perp_{AB}$ iff $\pi_{CA}; R; \pi_{BD} = \pi_{CD}$.



Important Constructions I

Definition

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For L -relations any singleton set is a unit.

In addition, for arrow categories $\pi_{11} = \mathbb{I}_1$ is sufficient since for every A we have

$$\begin{aligned}\pi_{A1}; \pi_{1A} &= \pi_{AA} \\ &\supseteq \mathbb{I}_A.\end{aligned}$$

\mathcal{A} uniform



Important Constructions II

Definition

An object $A \times B$ together with relations $\pi : A \times B \rightarrow A$ and $\rho : A \times B \rightarrow B$ is called a relational product of A and B iff

$$\pi, \rho \text{ are crisp, } \pi^\sim; \pi \sqsubseteq \mathbb{I}_A, \quad \rho^\sim; \rho \sqsubseteq \mathbb{I}_B, \quad \pi^\sim; \rho = \Pi_{AB}, \quad \pi; \pi^\sim \sqcap \rho; \rho^\sim = \mathbb{I}_{A \times B}.$$



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$$\pi = \begin{pmatrix} 1 & 0 \\ 1 & 0 \\ 1 & 0 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \end{pmatrix}, \quad \rho = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$



Important Constructions III

Definition

An object L^A together with a relation $\varepsilon : A \rightarrow L^A$ is called a fuzzy relational power iff

$$\text{syQ}(\varepsilon, \varepsilon)^\downarrow \sqsubseteq \mathbb{I}_{L^A} \quad \text{and} \quad \text{syQ}(R, \varepsilon)^\downarrow \text{ is total for every } R : A \rightarrow B.$$



Important Constructions III

Definition

An object L^A together with a relation $\varepsilon : A \rightarrow L^A$ is called a fuzzy relational power iff

$$\text{syQ}(\varepsilon, \varepsilon)^\downarrow \sqsubseteq \mathbb{I}_{L^A} \quad \text{and} \quad \text{syQ}(R, \varepsilon)^\downarrow \text{ is total for every } R : A \rightarrow B.$$

$$L = \begin{array}{c} 1 \\ | \\ a \\ | \\ 0 \end{array} \quad \varepsilon = \begin{pmatrix} 0 & 0 & 0 & a & a & a & 1 & 1 & 1 \\ 0 & a & 1 & 0 & a & 1 & 0 & a & 1 \end{pmatrix}$$

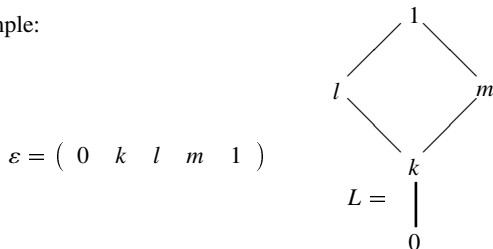
Notation: $\Lambda(R) = \text{syQ}(R^\sim, \varepsilon)^\downarrow : A \rightarrow L^B$ if $R : A \rightarrow B$.



Internal Version of L

The object L^1 acts as an internal version of the lattice L .

Example:



Every scalar $\alpha : 1 \rightarrow 1$ corresponds to a crisp map $p : 1 \rightarrow L$ (also called crisp point), i.e., $\Lambda(\alpha); \varepsilon^- = \alpha$ for every scalar $\alpha : 1 \rightarrow 1$ and $\Lambda(p; \varepsilon^-) = p$ for every crisp point $p : 1 \rightarrow L$.

The crisp order relation $(\varepsilon \setminus \varepsilon)^\downarrow$ corresponds to the order on scalars, i.e., $\alpha \sqsubseteq \beta$ iff $\text{syQ}(\beta, \varepsilon)^\downarrow \sqsubseteq \text{syQ}(\alpha, \varepsilon)^\downarrow; (\varepsilon \setminus \varepsilon)^\downarrow$.



Internal Version of L

The whole lattice structure as well as the arrow operations can be recovered on the interval version of L .

$$\text{zero} = \Lambda(\perp_{11}),$$

$$\text{one} = \Lambda(\mathbb{I}_1),$$

$$\mathcal{M}_2 = \Lambda(\pi; \varepsilon^\sim \sqcap \rho; \varepsilon^\sim),$$

$$\mathcal{J}_2 = \Lambda(\pi; \varepsilon^\sim \sqcup \rho; \varepsilon^\sim),$$

$$\text{down} = \Lambda(\varepsilon^\sim \downarrow),$$

$$\text{up} = \Lambda(\varepsilon^\sim \uparrow)$$



Covectorization

Given a relation $R : A \rightarrow B$ we may define a (co)vector $\text{cov}(R) : A \times B \rightarrow 1$ by:

$$\text{cov}(R) := (\pi; R \sqcap \rho); \Pi_{B1}$$

Given a (co)vector $v : A \times B \rightarrow 1$ we may define:

$$\text{rel}(v) := \pi^{\sim}; (v; \Pi_{1B} \sqcap \rho)$$

We have

$$\text{rel}(\text{cov}(R)) = R \text{ and } \text{cov}(\text{rel}(v)) = v.$$



Covectorization II

We define

$$\begin{aligned}\text{swap}_{AB} &:= \rho; \pi^{\sim} \sqcap \pi; \rho^{\sim}, \\ \text{comp}_{ABC} &:= (\mathbb{I}_{(A \times B) \times (B \times C)} \sqcap \pi; \rho; \pi^{\sim}; \rho^{\sim}); (\pi; \pi; \pi^{\sim} \otimes \rho; \rho; \rho^{\sim}).\end{aligned}$$

Then we have:

Lemma

Let be $Q, R : A \rightarrow B$, and $S : B \rightarrow C$. Then we have:

- 1 $\text{cov}(Q^{\downarrow}) = \text{cov}(Q)^{\downarrow}$,
- 2 $\text{cov}(Q^{\uparrow}) = \text{cov}(Q)^{\uparrow}$,
- 3 $\text{cov}(Q \sqcup R) = \text{cov}(Q) \sqcup \text{cov}(R)$,
- 4 $\text{cov}(Q \sqcap R) = \text{cov}(Q) \sqcap \text{cov}(R)$,
- 5 $\text{cov}(Q^{\sim}) = \text{swap}; \text{cov}(Q)$,
- 6 $\text{cov}(R; S) = \text{comp}^{\sim}; (\pi; \text{cov}(R) \sqcap \rho; \text{cov}(S))$.



Change of Base I

We start with two arrow categories \mathcal{A}_1 and \mathcal{A}_2 each with a unit, relational products, and relational powers over the Heyting algebras L_1 and L_2 , respectively, so that the Dedekind categories \mathcal{A}_1^\downarrow and \mathcal{A}_2^\downarrow are isomorphic with the same objects, i.e., we will require that there are functors $H : \mathcal{A}_1 \rightarrow \mathcal{A}_2$ and $H^{-1} : \mathcal{A}_2 \rightarrow \mathcal{A}_1$ that are the identity on objects, preserve all operations and constants of \mathcal{A}_1^\downarrow resp. \mathcal{A}_2^\downarrow , and are inverse to each other.

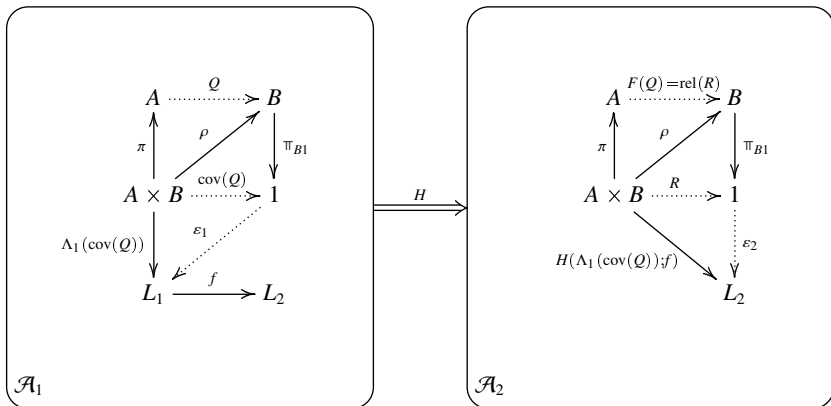
Given a crisp map $f : L_1 \rightarrow L_2$ in \mathcal{A}_1 , we define

$$F(Q) := \text{rel}(H(\Lambda_1(\text{cov}(Q))); f); \varepsilon_2^\sim)$$

for every $Q : A \rightarrow B$ in \mathcal{A}_1 .



Change of Base II



Change of Base III

Theorem

F preserves converse, i.e., $F(Q^-) = F(Q)^-$ for all relations $Q : A \rightarrow B$.



Change of Base IV

$$\mathbf{P1: } H((\mathbb{I}_{L_1} \sqcap \text{down}_1); f) = H(f); (\mathbb{I}_{L_2} \sqcap \text{down}_2),$$

$$\mathbf{P2a: } H(\text{one}_1; f) = \text{one}_2,$$

$$\mathbf{P2b: } \text{one}_2; H(f^-) = H(\text{one}_1).$$

Theorem

*If f satisfies **P1**, **P2a** and **P2b**, then F extends H , i.e., $F(Q) = H(Q)$ for all crisp relations $Q : A \rightarrow B$.*



Change of Base V

$$\mathbf{P3}: H(\mathcal{J}_2^1; f) = H(f \otimes f); \mathcal{J}_2^2.$$

Theorem

If f satisfies **P3**, then F preserves binary joins, i.e., $F(Q \sqcup R) = F(Q) \sqcup F(R)$ for all relations $Q, R : A \rightarrow B$.



Change of Base VI

$$\mathbf{P4}: H(\mathcal{M}_2^1; f) = H(f \otimes f); \mathcal{M}_2^2.$$

Theorem

If f satisfies **P4**, then F preserves binary meets, i.e., $F(Q \sqcap R) = F(Q) \sqcap F(R)$ for all relations $Q, R : A \rightarrow B$.



Change of Base VII

$$\mathbf{P5}: H(\Lambda_1(\varepsilon_1^{\downarrow}; \varepsilon_1^{\downarrow}); f) = \Lambda_2(H(\varepsilon_1^{\downarrow}; f)); \mathcal{J}^2.$$

Theorem

If f satisfies **P4** and **P5**, then F preserves composition, i.e., $F(R; S) = F(R); F(S)$ for all relations $R : A \rightarrow B$ and $S : B \rightarrow C$.



Thank you for your attention.

